

MIT 6.035 Spring 2011 Quiz 1 (100 points)

Your Full Name Here:

Your Athena ID Here:

1. (5 points) Write a regular expression for the language $L = \{0^n 1^m \mid (n + m) \text{ is even}\}$.

$$(00)^* (11)^* \mid 0(00)^* 1(11)^*$$

2. (20 points) Let the alphabet $\Sigma = \{0, 1\}$.

(a) (5 points) Write a regular expression for the language of all strings over Σ that contain the contiguous substring 11.

$$(0 \mid 1)^* 11 (0 \mid 1)^*$$

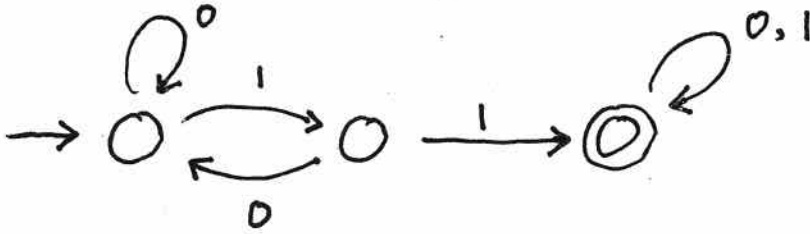
(b) (5 points) Write a regular expression for the language of all strings over Σ that do not contain the contiguous substring 11.

$$0^* \mid 0^* 1 (00^* 1)^* 0^*$$

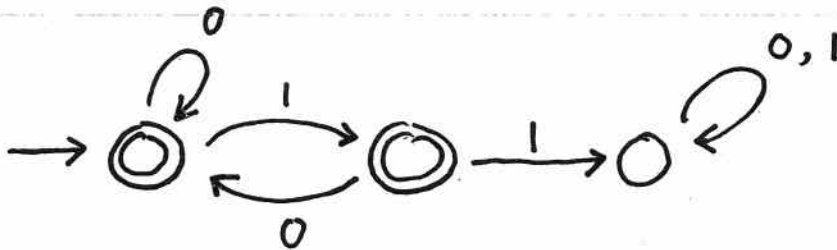
or

$$(10 \mid 0)^* (1 \mid \epsilon)$$

(c) (5 points) Give a non-deterministic finite automaton (NFA) for the language of all strings over Σ that contain the contiguous substring 11.

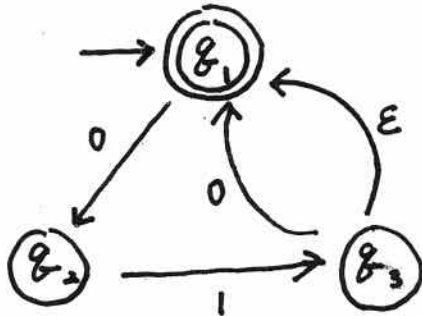


(d) (5 points) Give a non-deterministic finite automaton (NFA) for the language of all strings over Σ that don't contain the contiguous substring 11.

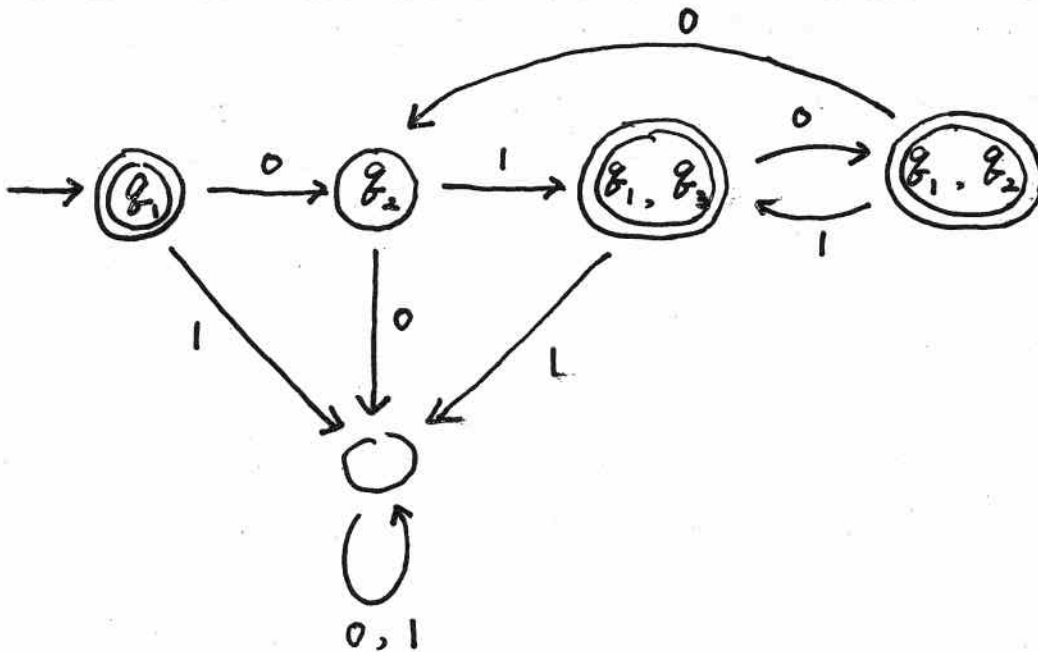


3. (10 points)

(a) (5 points) Give a non-deterministic finite automaton (NFA) for the language $L = (010 \mid 01)^*$. The NFA must contain at most 3 states. (Hint: draw an NFA with 4 states, then optimize).



(b) (5 points) Give a deterministic finite automaton for the language L.



4. (30 points)

Consider the following grammar:

$$S \rightarrow L = R$$
$$L \rightarrow *R \mid id$$
$$R \rightarrow L$$

You can think of L and R as standing for l-value and r-value, respectively. * is the dereference operator or indirection operator in C-like languages.

A shift-reduce parser can perform the following sequence of actions to accept the string “*id = id”.

shift » shift » reduce » reduce » reduce » shift » shift » reduce » reduce » reduce » accept

(a) (10 points) Give a sequence of actions that a shift-reduce parser can take to accept the string “id = id”.

shift → reduce → shift → shift → reduce → reduce
→ reduce → accept

(b) (10 points) Give a sequence of actions that a shift-reduce parser can take to accept the string “*id = *id”.

shift → shift → reduce → reduce → reduce
→ shift → shift → shift → reduce → reduce
→ reduce → reduce → reduce → accept

(c) (10 points) Is the grammar ambiguous? Why or why not?

No.

There doesn't exist a string which can be generated by the grammar in more than one way.

5. (15 points)

Consider the following grammar:

$S \rightarrow \text{if } E \text{ then } S \text{ else } S \mid \text{begin } S L \mid \text{print } E \mid \epsilon$

$L \rightarrow \text{end} \mid ; S L$

$E \rightarrow \text{num} = \text{num}$

The goal is to write a recursive-descent parser for the grammar. You are given the following $L()$ and $E()$ functions. Your job is to write the $S()$ function on the next page.

```
L() {
    if (token = end) {
        match(end);
    } else if (token = ;) {
        match(;); S(); L();
    } else {
        throw SyntaxError;
    }
}

E() {
    if (token = num) {
        match(num); match(=); match(num);
    } else {
        throw SyntaxError;
    }
}
```

```
S() {
```

```
    if (token = if) {  
        match(if); E(); match(then); S();  
        match(else); S();  
    } else if (token = begin) {  
        match(begin); S(); L();  
    } else if (token = print) {  
        match(print); E();  
    } else { // S → ε  
        // do nothing  
    }  
}
```

6. (20 points)

The following is a code snippet of legal-01.dcf:

```
class Program {
    int A[100];
    int length;

    void main() {
        int temp;

        length = 100;

        callout("srandom", 17);

        for i = 0, length {
            temp = callout("random");
            A[i] = temp;
        }

        /* <HERE> */
    }
}
```

What should the symbol tables look like at <HERE>, considering the semantics of the Decaf language? Complete the symbol tables on the next page in the similar way to the symbol tables presented at Lecture 5. (Hint: note that the Decaf language is different from the language presented at Lecture 5).

